Network Query Engines

Craig Knoblock
USC Information Sciences Institute
Overview

- Network Query Engines
  - Tukwila, Telegraph, Niagara
  - Dataflow & pipelining similar to Theseus
  - Execution system with support for efficient query execution from remote data sources
  - Automatically generate query plans from XML queries
  - No support for loops, conditionals, or external interactions
  - Designed for querying only, not monitoring (except for NiagaraCQ)
Tukwila (Ives et al. 1999)

- Adaptive network query processing for XML data
  - Interleaved execution and optimization
  - Inter-operator adaptivity
  - Dynamic operator re-ordering based on events
    - Memory overflow, wrapper timeout

- Notable new operators
  - X-SCAN: Efficient querying of streaming XML docs
  - JOIN: Double pipelined hash (probe is LHS or RHS)
  - DYNAMIC COLLECTOR: Efficient unioning of sources
Tukwila – Interleaved Planning and Execution

- Generates initial plan
- Can generate partial plans and expand them later
- Uses rules to decide when to re-optimize

From Ives et al., SIGMOD’99

WHEN end_of_fragment(0)
  IF card(result) > 100,000
  THEN re-optimize
**Tukwila – Adaptive Double Pipelined Hash Join**

From Ives et al., SIGMOD’99

### Hybrid Hash Join
- No output until inner read
- Asymmetric (inner vs. outer)

### Double Pipelined Hash Join
- Outputs data immediately
- Symmetric
- More memory

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Tukwila – Dynamic Collector Op

- Smart union operator
- Supports
  - Timeouts
  - slow sources
  - overlapping sources

From Ives et al., SIGMOD’99

WHEN
timeout(CustReviews)
DO activate(NYTimes),
activate(alt.books)
Niagara (Naughton, DeWitt, et al. 2000)

- Adaptive network query processing for XML data
  - Interleaved execution + document search
  - Supports streaming over blocking operators
    - Synchronization by re-evaluating operators or by propagating the differential result
Execution with partial results
[Shanmugasundaram et al. 2000]

- Niagara uses partial results to reduce the effects of blocking operators
  - Reduces blocking nature of aggregation or joins
- Basic idea
  - Execute future operators as data streams in, refine as slow operators catch up
  - Execution is driven by the availability of real data
  - Results are refined as additional data are processed
Approaches to Refining Results

• Re-evaluation
  • As new data becomes available, the operators re-output the results and the downstream operators are re-executed
  • Can be costly, but simple to implement

• Differential Algorithm
  • Each operator must support additions, deletes, and updates
  • Changed results must then be propagated to downstream operators
Telegraph (Hellerstein et al. 2000)

- Tuple-level adaptivity
- **Rivers** (optimize horizontal parallelism)
  - Adaptive dataflow on clusters (ie, data partitioning)
- **Eddies** (optimize vertical parallelism)
  - Leverage commutative property of query operators to dynamically route tuples for processing
Adaptable Joins, Issue 1

- **Synchronization Barriers**
  - One input frozen, waiting for the other
  - Can’t adapt while waiting for barrier!
  - So, favor joins that have:
    - no barriers
    - at worst, adaptable barriers
Adaptable Joins, Issue 2

- Would like to reorder *in-flight* (pipelined) joins
- Base case: swap inputs to a join
  - What about per-input state?

- Moment of symmetry:
  - inputs can be swapped w/o state management
- E.g.
  - Nested Loops: at the end of each inner loop
  - Merge Join: any time*
  - Hybrid or Grace Hash: never!

- More frequent moments of symmetry
  → more frequent adaptivity
**Ripple Joins: Prime for Adaptivity**

- **Ripple Joins**
  - Pipelined hash join (a.k.a. hash ripple, Xjoin)
    - No synchronization barriers
    - Continuous symmetry
    - Good for equi-join
  - Simple (or block) ripple join
    - Synchronization barriers at “corners”
    - Moments of symmetry at “corners”
    - Good for non-equi-join
  - Index nested loops
    - Short barriers
    - No symmetry
Beyond Binary Joins

- Think of swapping “inners”
  - Can be done at a global moment of symmetry
- Intuition: like an n-ary join
  - Except that each pair can be joined by a different algorithm!

- So...
  - Need to introduce n-ary joins to a traditional query engine

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Telegraph – Beyond Reordering Joins

Eddy

- A pipelining tuple-routing iterator (just like join or sort)
- Adjusts flow adaptively
  - Tuples flow in different orders
  - Visit each op once before output
- Naïve routing policy:
  - All ops fetch from eddy as fast as possible
  - Previously-seen tuples precede new tuples

From Avnur & Hellerstein, SIGMOD 2000
Discussion

• Theseus, Tukwila, Telegraph, Niagara are all:
  • Streaming dataflow systems
  • Targeting network-based query processing
    • Large source latencies
    • Unknown characteristics of sources
  • Proposed various techniques for improving the efficiency of processing data
    • More efficient operators (e.g., double-pipelined join)
    • Tuple-level adaptivity
    • Partial results for blocking operators
    • Speculative execution